Virtual memory | A storage allocation scheme in which secondary memory can be addressed as though it were part of main memory. The addresses a program may use to reference memory are distinguished from the addresses the memory system uses to identify physical storage sites, and program-generated addresses are translated automatically to the corresponding machine addresses. The size of virtual storage is limited by the addressing scheme of the computer system and by the amount of secondary memory available and not by the actual number of main storage locations.

Virtual address | The address assigned to a location in virtual memory to allow that location to be accessed as though it were part of main memory.

Virtual address space | The virtual storage assigned to a process.

Address space | The range of memory addresses available to a process.

Real address | The address of a storage location in main memory.

Table 8.1 Virtual Memory Terminology
Hardware and Control Structures

Two characteristics fundamental to memory management:

1) all memory references are logical addresses that are dynamically translated into physical addresses at run time
2) a process may be broken up into a number of pieces that don’t need to be contiguously located in main memory during execution

If these two characteristics are present, it is not necessary that all of the pages or segments of a process be in main memory during execution
Operating system brings into main memory a few pieces of the program

- Resident set
  - portion of process that is in main memory

- An interrupt is generated when an address is needed that is not in main memory

- Operating system places the process in a blocking state

Continued . . .
Execution of a Process

- Piece of process that contains the logical address is brought into main memory
- Operating system issues a disk I/O Read request
- Another process is dispatched to run while the disk I/O takes place
- An interrupt is issued when disk I/O is complete, which causes the operating system to place the affected process in the Ready state
Implications

- More processes may be maintained in main memory
  - only load in some of the pieces of each process
  - with so many processes in main memory, it is very likely a process will be in the Ready state at any particular time
- A process may be larger than all of main memory
Real and Virtual Memory

**Real memory**
- main memory, the actual RAM

**Virtual memory**
- memory on disk
- allows for effective multiprogramming and relieves the user of tight constraints of main memory
<table>
<thead>
<tr>
<th>Simple Paging</th>
<th>Virtual Memory Paging</th>
<th>Simple Segmentation</th>
<th>Virtual Memory Segmentation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Main memory partitioned into small fixed-size chunks called frames</td>
<td></td>
<td>Main memory not partitioned</td>
<td></td>
</tr>
<tr>
<td>Program broken into pages by the compiler or memory management system</td>
<td>Program segments specified by the programmer to the compiler (i.e., the decision is made by the programmer)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Internal fragmentation within frames</td>
<td>No internal fragmentation</td>
<td></td>
<td></td>
</tr>
<tr>
<td>No external fragmentation</td>
<td>External fragmentation</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Operating system must maintain a page table for each process showing which frame each page occupies</td>
<td>Operating system must maintain a segment table for each process showing the load address and length of each segment</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Operating system must maintain a free frame list</td>
<td>Operating system must maintain a list of free holes in main memory</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Processor uses page number, offset to calculate absolute address</td>
<td>Processor uses segment number, offset to calculate absolute address</td>
<td></td>
<td></td>
</tr>
<tr>
<td>All the pages of a process must be in main memory for process to run, unless overlays are used</td>
<td>Not all pages of a process need be in main memory frames for the process to run. Pages may be read in as needed</td>
<td>All the segments of a process must be in main memory for process to run, unless overlays are used</td>
<td>Not all segments of a process need be in main memory for the process to run. Segments may be read in as needed</td>
</tr>
<tr>
<td>Reading a page into main memory may require writing a page out to disk</td>
<td></td>
<td>Reading a segment into main memory may require writing one or more segments out to disk</td>
<td></td>
</tr>
</tbody>
</table>

**Table 8.2**

Characteristics of Paging and Segmentation
A state in which the system spends most of its time swapping process pieces rather than executing instructions.

To avoid this, the operating system tries to guess, based on recent history, which pieces are least likely to be used in the near future.
Principle of Locality

- Program and data references within a process tend to cluster.
- Only a few pieces of a process will be needed over a short period of time.
- Therefore it is possible to make intelligent guesses about which pieces will be needed in the future.
- Avoids thrashing.
For virtual memory to be practical and effective:

- hardware must support paging and segmentation
- operating system must include software for managing the movement of pages and/or segments between secondary memory and main memory
The term *virtual memory* is usually associated with systems that employ paging.

Use of paging to achieve virtual memory was first reported for the Atlas computer.

Each process has its own page table:
- each page table entry contains the frame number of the corresponding page in main memory.
(a) Paging only

Virtual Address

<table>
<thead>
<tr>
<th>Page Number</th>
<th>Offset</th>
</tr>
</thead>
</table>

Page Table Entry

<table>
<thead>
<tr>
<th>Page Number</th>
<th>Frame Number</th>
</tr>
</thead>
</table>

(b) Segmentation only

Virtual Address

<table>
<thead>
<tr>
<th>Segment Number</th>
<th>Offset</th>
</tr>
</thead>
</table>

Segment Table Entry

<table>
<thead>
<tr>
<th>Segment Number</th>
<th>Length</th>
<th>Segment Base</th>
</tr>
</thead>
</table>

(c) Combined segmentation and paging

Virtual Address

<table>
<thead>
<tr>
<th>Segment Number</th>
<th>Page Number</th>
<th>Offset</th>
</tr>
</thead>
</table>

Segment Table Entry

<table>
<thead>
<tr>
<th>Control Bits</th>
<th>Length</th>
<th>Segment Base</th>
</tr>
</thead>
</table>

Page Table Entry

<table>
<thead>
<tr>
<th>Page Number</th>
<th>Frame Number</th>
<th>Control Bits</th>
<th>Length</th>
<th>Segment Base</th>
</tr>
</thead>
</table>

P = present bit
M = Modified bit

Figure 8.1 Typical Memory Management Formats
Figure 8.2 Address Translation in a Paging System
Figure 8.3 A Two-Level Hierarchical Page Table
Figure 8.4 Address Translation in a Two-Level Paging System
Inverted Page Table

- Page number portion of a virtual address is mapped into a hash value
  - hash value points to inverted page table
- Fixed proportion of real memory is required for the tables regardless of the number of processes or virtual pages supported
- Structure is called *inverted* because it indexes page table entries by frame number rather than by virtual page number
Figure 8.5  Inverted Page Table Structure
Inverted Page Table

Each entry in the page table includes:

- **Page number**: the process that owns this page
- **Process identifier**: includes flags and protection and locking information
- **Control bits**: the index value of the next entry in the chain
- **Chain pointer**:
Translation Lookaside Buffer (TLB)

- Each virtual memory reference can cause two physical memory accesses:
  - one to fetch the page table entry
  - one to fetch the data

- To overcome the effect of doubling the memory access time, most virtual memory schemes make use of a special high-speed cache called a translation lookaside buffer
Figure 8.6 Use of a Translation Lookaside Buffer
Figure 8.7  Operation of Paging and Translation Lookaside Buffer (TLB) [FURH87]
Associative Mapping

- The TLB only contains some of the page table entries so we cannot simply index into the TLB based on page number.
- Each TLB entry must include the page number as well as the complete page table entry.
- The processor is equipped with hardware that allows it to interrogate simultaneously a number of TLB entries to determine if there is a match on page number.
Figure 8.8   Direct Versus Associative Lookup for Page Table Entries
Figure 8.9  Translation Lookaside Buffer and Cache Operation
The smaller the page size, the lesser the amount of internal fragmentation
- however, more pages are required per process
- more pages per process means larger page tables
- for large programs in a heavily multiprogrammed environment some portion of the page tables of active processes must be in virtual memory instead of main memory
- the physical characteristics of most secondary-memory devices favor a larger page size for more efficient block transfer of data
Figure 8.10  Typical Paging Behavior of a Program

(a) Page Size

\[ P = \text{size of entire process} \]
\[ W = \text{working set size} \]
\[ N = \text{total number of pages in process} \]

(b) Number of Page Frames Allocated

\[ \text{Page Fault Rate} \]

\[ P \]

\[ W \]

\[ N \]
<table>
<thead>
<tr>
<th>Computer</th>
<th>Page Size</th>
</tr>
</thead>
<tbody>
<tr>
<td>Atlas</td>
<td>512 48-bit words</td>
</tr>
<tr>
<td>Honeywell-Multics</td>
<td>1024 36-bit words</td>
</tr>
<tr>
<td>IBM 370/XA and 370/ESA</td>
<td>4 Kbytes</td>
</tr>
<tr>
<td>VAX family</td>
<td>512 bytes</td>
</tr>
<tr>
<td>IBM AS/400</td>
<td>512 bytes</td>
</tr>
<tr>
<td>DEC Alpha</td>
<td>8 Kbytes</td>
</tr>
<tr>
<td>MIPS</td>
<td>4 Kbytes to 16 Mbytes</td>
</tr>
<tr>
<td>UltraSPARC</td>
<td>8 Kbytes to 4 Mbytes</td>
</tr>
<tr>
<td>Pentium</td>
<td>4 Kbytes or 4 Mbytes</td>
</tr>
<tr>
<td>IBM POWER</td>
<td>4 Kbytes</td>
</tr>
<tr>
<td>Itanium</td>
<td>4 Kbytes to 256 Mbytes</td>
</tr>
</tbody>
</table>

**Table 8.3**

Example Page Sizes
Contemporary programming techniques used in large programs tend to decrease the locality of references within a process.

- The design issue of page size is related to the size of physical main memory and program size.

- Main memory is getting larger and address space used by applications is also growing.

- Most obvious on personal computers where applications are becoming increasingly complex.
Segmentation

Advantages:

- simplifies handling of growing data structures
- allows programs to be altered and recompiled independently
- lends itself to sharing data among processes
- lends itself to protection
Segment Organization

- Each segment table entry contains the starting address of the corresponding segment in main memory and the length of the segment.
- A bit is needed to determine if the segment is already in main memory.
- Another bit is needed to determine if the segment has been modified since it was loaded in main memory.
Figure 8.11  Address Translation in a Segmentation System
In a combined paging/segmentation system, a user’s address space is broken up into a number of segments. Each segment is broken up into a number of fixed-sized pages which are equal in length to a main memory frame.

Segmentation is visible to the programmer

Paging is transparent to the programmer
Figure 8.12  Address Translation in a Segmentation/Paging System
### Figure 8.1 Typical Memory Management Formats

#### (a) Paging only

<table>
<thead>
<tr>
<th>Virtual Address</th>
<th>Segment Number</th>
<th>Page Number</th>
<th>Offset</th>
</tr>
</thead>
</table>

#### (b) Segmentation only

| Segment Table Entry
<table>
<thead>
<tr>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>Control Bits</td>
</tr>
</tbody>
</table>

#### (c) Combined segmentation and paging

| Page Table Entry
<table>
<thead>
<tr>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>P</td>
</tr>
</tbody>
</table>

P = present bit
M = Modified bit

(c) Combined segmentation and paging
Protection and Sharing

- Segmentation lends itself to the implementation of protection and sharing policies.
- Each entry has a base address and length so inadvertent memory access can be controlled.
- Sharing can be achieved by segments referencing multiple processes.
Figure 8.13 Protection Relationships Between Segments
The design of the memory management portion of an operating system depends on three fundamental areas of choice:

- whether or not to use virtual memory techniques
- the use of paging or segmentation or both
- the algorithms employed for various aspects of memory management
<table>
<thead>
<tr>
<th>Fetch Policy</th>
<th>Resident Set Management</th>
</tr>
</thead>
<tbody>
<tr>
<td>Demand paging</td>
<td>Resident set size</td>
</tr>
<tr>
<td>Prepaging</td>
<td>Fixed</td>
</tr>
<tr>
<td></td>
<td>Variable</td>
</tr>
<tr>
<td>Placement Policy</td>
<td>Replacement Scope</td>
</tr>
<tr>
<td>Replacement Policy</td>
<td>Global</td>
</tr>
<tr>
<td>Basic Algorithms</td>
<td>Local</td>
</tr>
<tr>
<td>Optimal</td>
<td></td>
</tr>
<tr>
<td>Least recently used (LRU)</td>
<td></td>
</tr>
<tr>
<td>First-in-first-out (FIFO)</td>
<td></td>
</tr>
<tr>
<td>Clock</td>
<td></td>
</tr>
<tr>
<td>Page Buffering</td>
<td>Cleaning Policy</td>
</tr>
<tr>
<td></td>
<td>Demand</td>
</tr>
<tr>
<td></td>
<td>Precleaning</td>
</tr>
<tr>
<td>Load Control</td>
<td>Degree of multiprogramming</td>
</tr>
</tbody>
</table>

Table 8.4  Operating System Policies for Virtual Memory
Determine when a page should be brought into memory.

Two main types:

- Demand Paging
- Prepaging
Demand Paging

- Demand Paging
  - only brings pages into main memory when a reference is made to a location on the page
  - many page faults when process is first started
  - principle of locality suggests that as more and more pages are brought in, most future references will be to pages that have recently been brought in, and page faults should drop to a very low level
Prepaging

- Prepaging
  - pages other than the one demanded by a page fault are brought in
  - exploits the characteristics of most secondary memory devices
  - if pages of a process are stored contiguously in secondary memory it is more efficient to bring in a number of pages at one time
  - ineffective if extra pages are not referenced
  - should not be confused with “swapping”
Placement Policy

- Determines where in real memory a process piece is to reside
- Important design issue in a segmentation system
- Paging or combined paging with segmentation placing is irrelevant because hardware performs functions with equal efficiency
- For NUMA systems an automatic placement strategy is desirable
Replacement Policy

- Deals with the selection of a page in main memory to be replaced when a new page must be brought in.
  - Objective is that the page that is removed be the page least likely to be referenced in the near future.

- The more elaborate the replacement policy the greater the hardware and software overhead to implement it.
When a frame is locked the page currently stored in that frame may not be replaced

- kernel of the OS as well as key control structures are held in locked frames
- I/O buffers and time-critical areas may be locked into main memory frames
- locking is achieved by associating a lock bit with each frame
Algorithms used for the selection of a page to replace:

- Optimal
- Least recently used (LRU)
- First-in-first-out (FIFO)
- Clock
<table>
<thead>
<tr>
<th>Page address stream</th>
<th>2</th>
<th>3</th>
<th>2</th>
<th>1</th>
<th>5</th>
<th>2</th>
<th>4</th>
<th>5</th>
<th>3</th>
<th>2</th>
<th>5</th>
<th>2</th>
</tr>
</thead>
<tbody>
<tr>
<td>OPT</td>
<td>2</td>
<td>3</td>
<td>2</td>
<td>2</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
</tr>
<tr>
<td>LRU</td>
<td>2</td>
<td>3</td>
<td>3</td>
<td>2</td>
<td>5</td>
<td>5</td>
<td>4</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
<td>5</td>
</tr>
<tr>
<td>FIFO</td>
<td>2</td>
<td>3</td>
<td>3</td>
<td>2</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>CLOCK</td>
<td>2*</td>
<td>3*</td>
<td>2*</td>
<td>3*</td>
<td>5*</td>
<td>5*</td>
<td>5*</td>
<td>5*</td>
<td>5*</td>
<td>5*</td>
<td>5*</td>
<td>5*</td>
</tr>
</tbody>
</table>

F = page fault occurring after the frame allocation is initially filled

Figure 8.14 Behavior of Four Page-Replacement Algorithms
Least Recently Used (LRU)

- Replaces the page that has not been referenced for the longest time
- By the principle of locality, this should be the page least likely to be referenced in the near future
- Difficult to implement
  - one approach is to tag each page with the time of last reference
  - this requires a great deal of overhead
First-in-First-out (FIFO)

- Treats page frames allocated to a process as a circular buffer
- Pages are removed in round-robin style
  - simple replacement policy to implement
- Page that has been in memory the longest is replaced
Clock Policy

- Requires the association of an additional bit with each frame
  - referred to as the use bit
- When a page is first loaded in memory or referenced, the use bit is set to 1
- The set of frames is considered to be a circular buffer
- Any frame with a use bit of 1 is passed over by the algorithm
- Page frames visualized as laid out in a circle
Figure 8.15  Example of Clock Policy Operation
Figure 8.16 Comparison of Fixed-Allocation, Local Page Replacement Algorithms
Page Buffering

- Improves paging performance and allows the use of a simpler page replacement policy

A replaced page is not lost, but rather assigned to one of two lists

- Free page list
  - list of page frames available for reading in pages

- Modified page list
  - pages are written out in clusters
Replacement Policy and Cache Size

- With large caches, replacement of pages can have a performance impact
  - if the page frame selected for replacement is in the cache, that cache block is lost as well as the page that it holds
  - in systems using page buffering, cache performance can be improved with a policy for page placement in the page buffer
  - most operating systems place pages by selecting an arbitrary page frame from the page buffer
The OS must decide how many pages to bring into main memory:
- the smaller the amount of memory allocated to each process, the more processes can reside in memory.
- small number of pages loaded increases page faults.
- beyond a certain size, further allocations of pages will not effect the page fault rate.
Resident Set Size

**Fixed-allocation**
- Gives a process a fixed number of frames in main memory within which to execute
  - When a page fault occurs, one of the pages of that process must be replaced

**Variable-allocation**
- Allows the number of page frames allocated to a process to be varied over the lifetime of the process
The scope of a replacement strategy can be categorized as **global** or **local**

- **Global**
  - considers all unlocked pages in main memory

- **Local**
  - chooses only among the resident pages of the process that generated the page fault
<table>
<thead>
<tr>
<th>Fixed Allocation</th>
<th>Local Replacement</th>
<th>Global Replacement</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>• Number of frames allocated to a process is fixed.</td>
<td>• Not possible.</td>
</tr>
<tr>
<td></td>
<td>• Page to be replaced is chosen from among the frames allocated to that process.</td>
<td></td>
</tr>
<tr>
<td>Variable Allocation</td>
<td>• The number of frames allocated to a process may be changed from time to time to maintain the working set of the process.</td>
<td>• Page to be replaced is chosen from all available frames in main memory; this causes the size of the resident set of processes to vary.</td>
</tr>
<tr>
<td></td>
<td>• Page to be replaced is chosen from among the frames allocated to that process.</td>
<td></td>
</tr>
</tbody>
</table>

Table 8.5  Resident Set Management
**Fixed Allocation, Local Scope**

- Necessary to decide ahead of time the amount of allocation to give a process

- If allocation is too small, there will be a high page fault rate

  - increased processor idle time
  - increased time spent in swapping

- If allocation is too large, there will be too few programs in main memory
Variable Allocation

Global Scope

- Easiest to implement
  - adopted in a number of operating systems

- OS maintains a list of free frames

- Free frame is added to resident set of process when a page fault occurs

- If no frames are available the OS must choose a page currently in memory

- One way to counter potential problems is to use page buffering
When a new process is loaded into main memory, allocate to it a certain number of page frames as its resident set.

When a page fault occurs, select the page to replace from among the resident set of the process that suffers the fault.

Reevaluate the allocation provided to the process and increase or decrease it to improve overall performance.
Decision to increase or decrease a resident set size is based on the assessment of the likely future demands of active processes.

Key elements:

- criteria used to determine resident set size
- the timing of changes
<table>
<thead>
<tr>
<th>Sequence of Page References</th>
<th>Window Size,</th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>24</td>
<td>2</td>
<td>3</td>
<td>4</td>
<td>5</td>
</tr>
<tr>
<td>15</td>
<td>24</td>
<td>24</td>
<td>24</td>
<td>24</td>
</tr>
<tr>
<td>18</td>
<td>24 15</td>
<td>24 15</td>
<td>24 15</td>
<td>24 15</td>
</tr>
<tr>
<td>18</td>
<td>24 15 18</td>
<td>24 15 18</td>
<td>24 15 18</td>
<td>24 15 18</td>
</tr>
<tr>
<td>23</td>
<td>24 15 18 23</td>
<td>24 15 18 23</td>
<td>24 15 18 23</td>
<td>24 15 18 23</td>
</tr>
<tr>
<td>24</td>
<td>18 23 24</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>17</td>
<td>24 17 23 24</td>
<td>18 23 24 17</td>
<td>15 18 23 24</td>
<td>17 18 23 24</td>
</tr>
<tr>
<td>18</td>
<td>17 18 24 17</td>
<td>•</td>
<td>18 23 24 17</td>
<td>•</td>
</tr>
<tr>
<td>24</td>
<td>18 24 24 17</td>
<td>•</td>
<td>24 17 18</td>
<td>•</td>
</tr>
<tr>
<td>18</td>
<td>• 24 18 17</td>
<td>• 24 17 18</td>
<td>24 17 18 24</td>
<td></td>
</tr>
<tr>
<td>17</td>
<td>18 17 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
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<tr>
<td>17</td>
<td>17 18 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>17</td>
<td>17 18 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>17</td>
<td>18 17 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>17</td>
<td>17 18 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>17</td>
<td>18 17 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>18</td>
<td>24 17 24 18</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
</tbody>
</table>

**Figure 8.17 Working Set of Process as Defined by Window Size**
Figure 8.18  Typical Graph of Working Set Size [MAEK87]
Page Fault Frequency (PFF)

- Requires a use bit to be associated with each page in memory
- Bit is set to 1 when that page is accessed
- When a page fault occurs, the OS notes the virtual time since the last page fault for that process
- Does not perform well during the transient periods when there is a shift to a new locality
Evaluates the working set of a process at sampling instances based on elapsed virtual time

Driven by three parameters:

- The minimum duration of the sampling interval
- The maximum duration of the sampling interval
- The number of page faults that are allowed to occur between sampling instances
Cleaning Policy

- Concerned with determining when a modified page should be written out to secondary memory

Demand Cleaning
- A page is written out to secondary memory only when it has been selected for replacement

Precleaning
- Allows the writing of pages in batches
Load Control

- Determines the number of processes that will be resident in main memory
  - multiprogramming level
- Critical in effective memory management
- Too few processes, many occasions when all processes will be blocked and much time will be spent in swapping
- Too many processes will lead to thrashing
Figure 8.19 Multiprogramming Effects

Processor Utilization vs. Multiprogramming Level

Figure 8.19 Multiprogramming Effects
If the degree of multiprogramming is to be reduced, one or more of the currently resident processes must be swapped out.

Six possibilities exist:

- lowest-priority process
- faulting process
- last process activated
- process with the smallest resident set
- largest process
- process with the largest remaining execution window
UNIX

- Intended to be machine independent so its memory management schemes will vary
- Early UNIX: variable partitioning with no virtual memory scheme
- Current implementations of UNIX and Solaris make use of

SVR4 and Solaris use two separate schemes:

- paging system
- kernel memory allocator
Paging System and Kernel Memory Allocator

**Paging System**
- Provides a virtual memory capability that allocates page frames in main memory to processes.
- Allocates page frames to disk block buffers.

**Kernel Memory Allocator**
- Allocates memory for the kernel.
### UNIX SVR4 Memory Management Formats

#### (a) Page table entry

<table>
<thead>
<tr>
<th>Page frame number</th>
<th>Age</th>
<th>Copy on write</th>
<th>Modify</th>
<th>Reference</th>
<th>Valid</th>
<th>Protect</th>
</tr>
</thead>
</table>

#### (b) Disk block descriptor

<table>
<thead>
<tr>
<th>Swap device number</th>
<th>Device block number</th>
<th>Type of storage</th>
</tr>
</thead>
</table>

#### (c) Page frame data table entry

<table>
<thead>
<tr>
<th>Page state</th>
<th>Reference count</th>
<th>Logical device</th>
<th>Block number</th>
<th>Pfdata pointer</th>
</tr>
</thead>
</table>

#### (d) Swap-use table entry

<table>
<thead>
<tr>
<th>Reference count</th>
<th>Page/storage unit number</th>
</tr>
</thead>
</table>

**Figure 8.20** UNIX SVR4 Memory Management Formats
<table>
<thead>
<tr>
<th><strong>Page Table Entry</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Page frame number</strong></td>
</tr>
<tr>
<td>Refers to frame in real memory.</td>
</tr>
<tr>
<td><strong>Age</strong></td>
</tr>
<tr>
<td>Indicates how long the page has been in memory without being referenced. The length and contents of this field are processor dependent.</td>
</tr>
<tr>
<td><strong>Copy on write</strong></td>
</tr>
<tr>
<td>Set when more than one process shares a page. If one of the processes writes into the page, a separate copy of the page must first be made for all other processes that share the page. This feature allows the copy operation to be deferred until necessary and avoided in cases where it turns out not to be necessary.</td>
</tr>
<tr>
<td><strong>Modify</strong></td>
</tr>
<tr>
<td>Indicates page has been modified.</td>
</tr>
<tr>
<td><strong>Reference</strong></td>
</tr>
<tr>
<td>Indicates page has been referenced. This bit is set to 0 when the page is first loaded and may be periodically reset by the page replacement algorithm.</td>
</tr>
<tr>
<td><strong>Valid</strong></td>
</tr>
<tr>
<td>Indicates page is in main memory.</td>
</tr>
<tr>
<td><strong>Protect</strong></td>
</tr>
<tr>
<td>Indicates whether write operation is allowed.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th><strong>Disk Block Descriptor</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Swap device number</strong></td>
</tr>
<tr>
<td>Logical device number of the secondary device that holds the corresponding page. This allows more than one device to be used for swapping.</td>
</tr>
<tr>
<td><strong>Device block number</strong></td>
</tr>
<tr>
<td>Block location of page on swap device.</td>
</tr>
<tr>
<td><strong>Type of storage</strong></td>
</tr>
<tr>
<td>Storage may be swap unit or executable file. In the latter case, there is an indication as to whether the virtual memory to be allocated should be cleared first.</td>
</tr>
</tbody>
</table>

**Table 8.6**

**UNIX SVR4 Memory Management Parameters**

(Table can be found on page 379 in the textbook)
### Page Frame Data Table Entry

**Page state**
Indicates whether this frame is available or has an associated page. In the latter case, the status of the page is specified: on swap device, in executable file, or DMA in progress.

**Reference count**
Number of processes that reference the page.

**Logical device**
Logical device that contains a copy of the page.

**Block number**
Block location of the page copy on the logical device.

**Pfddata pointer**
Pointer to other pfddata table entries on a list of free pages and on a hash queue of pages.

### Swap-Use Table Entry

**Reference count**
Number of page table entries that point to a page on the swap device.

**Page/storage unit number**
Page identifier on storage unit.

---

**Table 8.6**

**UNIX SVR4 Memory Management Parameters**

(Table can be found on page 379 in the textbook)
Page Replacement

- The page frame data table is used for page replacement.
- Pointers are used to create lists within the table.
  - All available frames are linked together in a list of free frames available for bringing in pages.
  - When the number of available frames drops below a certain threshold, the kernel will steal a number of frames to compensate.
Figure 8.21  Two-Handed Clock Page-Replacement Algorithm
Kernel Memory Allocator

- The kernel generates and destroys small tables and buffers frequently during the course of execution, each of which requires dynamic memory allocation.

- Most of these blocks are significantly smaller than typical pages (therefore paging would be inefficient)

- Allocations and free operations must be made as fast as possible
Technique adopted for SVR4

UNIX often exhibits steady-state behavior in kernel memory demand
  i.e. the amount of demand for blocks of a particular size varies slowly in time

Defers coalescing until it seems likely that it is needed, and then coalesces as many blocks as possible
Initial value of $D_i$ is 0
After an operation, the value of $D_i$ is updated as follows

(I) if the next operation is a block allocate request:
   if there is any free block, select one to allocate
      if the selected block is locally free
         then $D_i := D_i + 2$
         else $D_i := D_i + 1$
   otherwise
      first get two blocks by splitting a larger one into two (recursive operation)
      allocate one and mark the other locally free
      $D_i$ remains unchanged (but D may change for other block sizes because of the recursive call)

(II) if the next operation is a block free request
  Case $D_i \geq 2$
     mark it locally free and free it locally
     $D_i := D_i - 2$
  Case $D_i = 1$
     mark it globally free and free it globally; coalesce if possible
     $D_i := 0$
  Case $D_i = 0$
     mark it globally free and free it globally; coalesce if possible
     select one locally free block of size $2^i$ and free it globally; coalesce if possible
     $D_i := 0$

Figure 8.22 Lazy Buddy System Algorithm
Linux

Memory Management

- Shares many characteristics with UNIX
- Is quite complex

Two main aspects

• process virtual memory
• kernel memory allocation
Three level page table structure:

- **Page directory**
  - process has a single page directory
  - each entry points to one page of the page middle directory
  - must be in main memory for an active process

- **Page middle directory**
  - may span multiple pages
  - each entry points to one page in the page table

- **Page table**
  - may also span multiple pages
  - each entry refers to one virtual page of the process
Figure 8.23  Address Translation in Linux Virtual Memory Scheme
Based on the clock algorithm

The use bit is replaced with an 8-bit age variable
  - incremented each time the page is accessed

Periodically decrements the age bits
  - a page with an age of 0 is an “old” page that has not been referenced in some time and is the best candidate for replacement

A form of least frequently used policy
Figure 8.24 Linux Page Reclaiming
Kernel memory capability manages physical main memory page frames
- primary function is to allocate and deallocate frames for particular uses

A buddy algorithm is used so that memory for the kernel can be allocated and deallocated in units of one or more pages
- Page allocator alone would be inefficient because the kernel requires small short-term memory chunks in odd sizes
- Slab allocation
  - used by Linux to accommodate small chunks

Possible owners of a frame include:
- user-space processes
- dynamically allocated kernel data
- static kernel code
- page cache
Virtual memory manager controls how memory is allocated and how paging is performed.

- Designed to operate over a variety of platforms.
- Uses page sizes ranging from 4 Kbytes to 64 Kbytes.
Windows Virtual Address Map

- On 32 bit platforms each user process sees a separate 32 bit address space allowing 4 Gbytes of virtual memory per process
  - by default half is reserved for the OS
- Large memory intensive applications run more effectively using 64-bit Windows
- Most modern PCs use the AMD64 processor architecture which is capable of running as either a 32-bit or 64-bit system
64-Kbyte region for NULL-pointer assignments (inaccessible)

64-Kbyte region for bad pointer assignments (inaccessible)

2-Gbyte user address space (unreserved, usable)

2-Gbyte region for the operating system (inaccessible)

Figure 8.25 Windows Default 32-bit Virtual Address Space
Windows Paging

- On creation, a process can make use of the entire user space of almost 2 Gbytes
- This space is divided into fixed-size pages managed in contiguous regions allocated on 64 Kbyte boundaries
- Regions may be in one of three states:
  - available
  - reserved
  - committed
Windows uses variable allocation, local scope

When activated, a process is assigned a data structure to manage its working set

Working sets of active processes are adjusted depending on the availability of main memory
Android includes a number of extensions to the normal Linux kernel memory management facility.

These include:

- **ASHMem**
  - This feature provides anonymous shared memory, which abstracts memory as file descriptors.
  - A file descriptor can be passed to another process to share memory.

- **Pmem**
  - This feature allocates virtual memory so that it is physically contiguous.
  - This feature is useful for hardware that does not support virtual memory.

- **Low Memory Killer**
  - This feature enables the system to notify an app or apps that they need to free up memory.
  - If an app does not cooperate, it is terminated.
Summary

- Hardware and control structures
  - Locality and virtual memory
  - Paging
  - Segmentation
  - Combined paging and segmentation
  - Protection and sharing

- OS software
  - Fetch policy
  - Placement policy
  - Replacement policy
  - Resident set management
  - Cleaning policy
  - Load control

- UNIX and Solaris memory management
  - Paging system
  - Kernel memory allocator

- Linux memory management
  - Linux virtual memory
  - Kernel memory allocation

- Windows memory management
  - Windows virtual address map
  - Windows paging
  - Windows 8 swapping

- Android memory management